

*Public Key Cryptography (Spring 2003)*

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**Note: Solving polynomial equations modulo prime powers**

We have seen Berlekamp's algorithm for finding square roots modulo a prime. However, this algorithm cannot be used directly for finding roots modulo a prime power  $p^e$  ( $e > 1$ ) because  $a^{\frac{p^e-1}{2}} = 1$  only holds for a small fraction of  $a \in \mathbb{Z}_{p^e}$ . We thus proceed as follows.

**Finding square roots modulo  $p^2$**

Given a prime  $p$  and  $a \in \mathbb{Z}_{p^2}$  we wish to find the solutions of

$$x^2 \equiv a \pmod{p^2} \tag{1}$$

We begin by solving the following (e.g., using Berlekamp's algorithm):

$$y^2 \equiv a \pmod{p} \tag{2}$$

Each solution  $x$  of (1) is congruent modulo  $p$  to a solution  $y$  of (2). It thus suffices to find, for each  $y$ , the values  $x = y + ip$  that fulfill (1). This is done as follows.

$$a \equiv (y + ip)^2 \equiv y^2 + 2yip + i^2p^2 \equiv y^2 + 2yip \pmod{p^2} \tag{3}$$

so we wish to say:

$$i \equiv \frac{a - y^2}{2yp} \pmod{p^2} \tag{4}$$

In general we cannot divide by a multiple of  $p$ , because it does not have an inverse modulo  $p^2$ . However, by (2) we have  $a - y^2 \equiv 0 \pmod{p}$ , so the  $p$  factors cancel out:  $a - y^2 \equiv tp \pmod{p^2}$  for some  $t$ , so we can compute  $i$ :

$$i = t/2y \pmod{p^2}$$

This fails if  $2y$  does not have an inverse modulo  $p^2$ , or equivalently, when  $\gcd(2y, p^2) > 1$ . This happens when either  $p = 2$  (which is an easy special case) or  $y$  is a multiple of  $p$ . In the latter case, if there exists a corresponding solution  $x = y + ip$  then  $a \equiv x^2 \equiv 0 \pmod{p^2}$ , so we merely need to consider the trivial case  $a = 0$  separately.

To conclude, we saw that each root modulo  $p$  can be "lifted" into a root modulo  $p^2$  by simple computation. This is a special case of "Hensel lifting".

### **Finding square roots modulo $p^e$ for $e > 2$**

To find roots modulo  $p^4$ , we first find roots modulo  $p$  and “lift” them into roots modulo  $p^2$  as above. Then, we then “lift” the roots modulo  $p^2$  to into roots modulo  $p^4$ . This can be similarly to the above, except we replace  $p$  by  $p^2$ , and  $p^2$  by  $p^4$ . The handling of non-invertible denominators can be generalized.

By such repeated squaring, we can compute roots modulo  $p^e$  for any  $e$  which is a power of 2. To compute roots modulo  $p^e$  for other  $e$ , simply compute the roots modulo  $p^{e'}$  where  $e' \geq e$  is a power of 2, and reduce them modulo  $p^e$ .

### **Finding roots of polynomials modulo $p^e$**

The above generalizes from the special case of solving  $x^2 = a$  (i.e., extracting square roots) to finding roots of arbitrary polynomials. The essential points are that in (3)  $f(y + ip) \pmod{p_i}$  is linear for any polynomial  $f$ , and that the cancellation of  $p$  in (4) always occurs.

### **Finding roots modulo arbitrary integers**

To compute roots modulo an arbitrary natural number  $n$  whose prime factorization is known to be  $n = p_1^{e_1} p_2^{e_2} \cdots p_l^{e_l}$ , first compute the roots modulo each of the  $p_i^{e_i}$  ( $i = 1, \dots, l$ ) and then combine them using the Chinese remainder theorem.